

100% Code Coverage is not equal 100% Code Coverage

How reliable is Code Coverage information shown in source code?

Code coverage is a method to assess the quality of test runs. Based on this evaluation new test cases are deduced, redundant test cases are eliminated and inefficient test cases are changed or replaced. At the same time, code coverage is a measure to find code that is never executed. This “dead code” represents undesired overhead when assigning memory resources and is also a potential hazard if this code was not executed during testing.

Code coverage measurement (mainly on source code level) is recommended for certification in avionic, medical, automotive and nuclear standards like DO-178B, DO248D, IEC 62304, ISO 26262 and SC45A.

According to Beizer [1] **100% statement coverage at source code level** might cover **75% or even fewer statements at object code level** [2].

In this article we will discuss where this gap comes from.

Basics

Programs written in high level programming languages like C, C++, ADA etc. are not running on a processor or controller without being converted to object code. It is important to keep the following in mind: starting from the source code, there is a chain of transformations that produces the object code running on the actual target. For C programs usually these steps are: preprocessing, compilation, optimization, assembly, linking and conversion. Details vary between systems (compilers, assemblers, linkers, converters). These steps are not bijective: information is added and removed in every step. For example, linking generates addresses for variables and throws away symbolic names.

These modifications directly affect code coverage analysis on source code level because

- often compiler or library calls add object code and
- it is not easy to detect when object code is added by a compiler or a library call and
- the added object code may not be 100% covered by your tests and
- it is very complex to establish which object code correlates to which source code line (usually code coverage analysis is shown in source code).

Reasons for object code coverage

The following examples show that every time you use *complex instructions* (e.g. a for loop), *complex conditions* (e.g. if statement) or *library/operating system functions* object code is added that is not directly traceable to source code statements. Sometimes also the *compiler generates object code*. So a lot of object code is not directly traceable to source code.

Consequently DO-178/ED-12B Section 6.4.4.2b states that if "... *the compiler generates object code that is not directly traceable to Source Code statements. Then, additional verification should be performed on the object code...*"

So if an application has to **conform to standards** like above then the **object code has to be covered too**. To get an application approved following three steps must be addressed (cite from CAST-12 [3]):

- *detect and identify any added object code*
- *establish the functionality of that code; and*
- *verify the correctness of the added, untraceable code sequences*

You can easily achieve object code verification and coverage using iSYSTEM on-chip debug tools and analyzers. iSYSTEM tools measure code coverage by recording all addresses being executed *on the processor* – measurement is done on low level binary data. No code instrumentation is needed. Code coverage results are displayed on source code, object code and assembly level and are correlations of the low level binary data to the high level source/object/assembly code. Object code insertions are easily visible in iSYSTEM's code coverage analysis window.

Prerequisite of such an object level code coverage measurement: Analysis on object code runs on the code loaded in the real target hardware. To measure code coverage a hardware tool connected to the target records the activities on the CPU buses (commands and data, external signals including time stamps). With this information it can be proved in object code whether code has been executed or not. This functionality is called tracing. Tracing is available through an In-Circuit-Emulator and also through an On-Chip Debugger (OCD), but the OCD interface has to have at least one code trace port to reconstruct and measure code execution. Perfer the usage of microcontrollers that implement on chip debug cells like NEXUS or ETM.

Relation between Source and Object Code

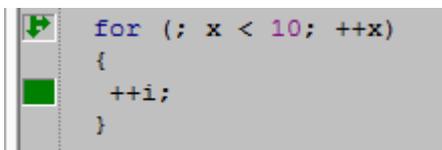
When reading the following examples, please note: **Non-conditional instructions** are marked with a **colored box**. This box is **green** for *executed* instructions and **red** for *non-executed* instructions. Red and green **arrows** display the **object level condition coverage** information for **conditional branches and conditional instructions**. The arrow is colored according to the paths that were executed.

1. Complex Instructions

The object code may not correspond to the source code as seen from the perspective of a source level programmer. Take this example:

```
1 int i,x;
2 for (; x < 10; ++x)
3 {
4   ++i;
5 }
```

Line 2 is typically (no matter of the processor architecture used) splitted in two or three blocks: one to initialize the loop, another to check the condition and increment the loop variable, and a tail to cleanup everything after it exits. The instructions are not run through sequentially but in two or three blocks that may be located in non-contiguous memory locations. Further, you have one or more conditions (see assembler below). Every condition splits the execution path to (at least) two possible execution paths. Code coverage information for all three blocks has to be merged to one source code line for visualization to the user. iSYSTEM tools display code coverage results depending on the executed code on the processor. In the disassembly window you can see the instruction blocks described above. Note the branch instruction “b” at the very beginning and the conditional branch “bc” with jump at the end.



```
for (; x < 10; ++x)
{
  ++i;
}
```

➔ 40008260:40008260	4800001C	for (; x < 10; ++x)	
➔ 40008264:40008264	801F000C	b	4000827C
➔ 40008268:40008268	30000001	++i;	
➔ 4000826C:4000826C	901F000C		
➔ 40008270:40008270	801F0008	for (; x < 10; ++x)	
➔ 40008274:40008274	30000001	lwz	r0,08(r31)
➔ 40008278:40008278	901F0008	addic	r0,r0,01
➔ 4000827C:4000827C	801F0008	stw	r0,08(r31)
➔ 40008280:40008280	2F800009	lwz	r0,08(r31)
➔ 40008284:40008284	409DFFE0	cmpi	7,0,r0,09
		bc	04,10,"test.c"::58 (40008264)



!! When measuring code coverage watch out that all paths are covered.

2. Complex Conditions

For example:

```

1 if (a == 3 || b ==4)
2 {
3   somefunction(c);
4 }

```

Line 1 of the example can be only partially tested when the first condition evaluates to true and the evaluation short-circuit logic* applies. In this case the second condition is not tested at all.

!! When measuring code coverage ensure that every condition is assessed.

iSYSTEM tools show this situation with the greatest possible amount of the detail. If the first condition is reached and the evaluation short-circuit logic applies, then the coverage information (in the source code window) will show the line as not executed. If you drill down on the assembler information for this line you will see that only part of the code was executed.

This is shown when using iSYSTEM tools: Note how the source code shows the "if" statement as executed (the arrowroot is green) but neither path taken (arrowheads are red). This is to signal a mixed case that has to be analyzed on the assembler level. In the disassembly window you see that the first condition was true, so the second one was never evaluated - therefore the coverage legend on these assembler lines shows red squares and arrows.

```

int a=3, b=4;

if (a==3 || b==4)
{
    somefunction();
}

```

4000828C:4000828C	801F000C	if (a==3 b==4)
40008290:40008290	2F800003	lwz r0,0C(r31)
40008294:40008294	419E0010	cmpi 7,0,r0,03
40008298:40008298	801F0008	bc 0C,1E,"test.c"::65 (400082A4)
4000829C:4000829C	2F800004	lwz r0,08(r31)
400082A0:400082A0	409E0008	cmpi 7,0,r0,04
		bc 04,1E,"test.c"::69 (400082A8)
		somefunction();
400082A4:400082A4	4BFFFF3D	bl somefunction (400081E0)

*Short-circuit evaluation, minimal evaluation, or McCarthy evaluation denotes the semantics of some Boolean operators in some programming languages in which the second argument is only executed or evaluated if the first argument does not suffice to determine the value of the expression: when the first argument of the AND function evaluates to false, the overall value must be false; and when the first argument of the OR function evaluates to true, the overall value must be true.

3. Libraries

Using library functions, e.g. operating system functions, will add object code where no source code is available. These library functions have to be thoroughly tested because they could include dead code or never taken code paths.

iSYSTEM tools gather coverage information directly from the hardware. So the analysis of library code isn't different to analysis from code written by the user. If source code is not available (e.g. for an operating system function) the information is given on the object and assembler levels only. iSYSTEM tools allow to configure if library code is included in the overall coverage information and reports or not.

This is shown when using iSYSTEM tools: Note that not all instructions of the library function were tested.

```

      _fixdfsi
20001F24: 7C0802A6  mflr      r0
20001F28: 9421FFD0  stwu     r1,-30(r1)
20001F2C: 90610008  stw      r3,08(r1)
20001F30: 9081000C  stw      r4,0C(r1)
20001F34: 38610008  la       r3,08(r1)
20001F38: 38810010  la       r4,10(r1)
20001F3C: 90010034  stw      r0,34(r1)
20001F40: 48000625  bl      _unpack_d (20002564)
20001F44: 80010010  lwz     r0,10(r1)
20001F48: 2F800002  cmpi    7,0,r0,02
20001F4C: 419E0030  bc      0C,1E,20001F7C
20001F50: 28800001  cmpli   7,0,r0,01
20001F54: 409D0028  bc      04,1D,20001F7C
20001F58: 2F800004  cmpi    7,0,r0,04
20001F5C: 409E0034  bc      04,1E,20001F90
20001F60: 80010014  lwz     r0,14(r1)
20001F64: 3C607FFF  lis     r3,7FFF0000
20001F68: 2F800000  cmpi    7,0,r0,00
20001F6C: 6063FFFF  ori     r3,FFFF
20001F70: 41BE0010  bc      0D,1E,20001F80
20001F74: 3C608000  lis     r3,80000000
20001F78: 48000008  b      20001F80
20001F7C: 38600000  li     r3,00
20001F80: 80010034  lwz     r0,34(r1)
20001F84: 38210030  addi   r1,30
20001F88: 7C0803A6  mtlr   r0

      _fixdfsi_EXIT_
20001F8C: 4E800020  blr
20001F90: 80010018  lwz     r0,18(r1)
20001F94: 2F800000  cmpi    7,0,r0,00
20001F98: 41BCFFE4  bc      0D,1C,20001F7C
20001F9C: 2F80001E  cmpi    7,0,r0,1E
20001FA0: 419DFFC0  bc      0C,1D,20001F60
20001FA4: 2100003C  subfic  r8,r0,3C
20001FA8: 3528FFE0  addic.  r9,r8,-20
20001FAC: 41800024  bc      0C,00,20001FD0
20001FB0: 80010020  lwz     r0,20(r1)
20001FB4: 7C044C30  srw     r4,r0,r9
20001FB8: 80010014  lwz     r0,14(r1)
20001FBC: 7C832378  mr      r3,r4
20001FC0: 2F800000  cmpi    7,0,r0,00
20001FC4: 41BEFFBC  bc      0D,1E,20001F80
20001FC8: 7C640000  neg     r3,r4
20001FCC: 4BFFFFB4  b      20001F80
20001FD0: 81410020  lwz     r10,20(r1)
20001FD4: 81210024  lwz     r9,24(r1)
20001FD8: 554B003C  rlwinm  r11,r10,1,0,30
20001FDC: 2008001F  subfic  r0,r8,1F
20001FE0: 7D6B0030  slw     r11,r11,r0
20001FE4: 7D244430  srw     r4,r9,r8
20001FE8: 7D642378  or      r4,r11,r4
20001FEC: 4BFFFFCC  b      20001FB8

```

Here you see iSYSTEM winIDEA's hierarchical code coverage view for above library function. Because there is no source code available for these functions the "Lines" information columns are empty. Nevertheless you can view full code coverage information for each function, each object code and each assembler instruction, or you can collapse parts of the tree to have an overview.

StatPane	Lines Bar	Lines	Sizes Bar	Sizes	Branches Bar	Branches
adddf3				0x78 / 0x78 (100%)		
addsf3				0x70 / 0x70 (100%)		
extendsfd2				0x350 / 0x3A0 (91%)		(26ne + 0t + 0nt + 0b) / 26 (100%)
fixdfsi				0x48 / 0xCC (35%)		(1ne + 2t + 5nt + 0b) / 8 (56%)
mflr r0				0x0 / 0x4 (0%)		
stwu r1, 30(r1)				0x0 / 0x4 (0%)		
stw r3, 08(r1)				0x0 / 0x4 (0%)		
20001F2C:20001F2C - 20001F2F:20001F2F				0x0 / 0x4 (0%)		
stw r4, 0C(r1)				0x0 / 0x4 (0%)		

4. Compiler Generated Code

Specially on smaller processors (e.g. Freescale MPC5604B), but not only, the compiler needs to introduce code to bridge the gap between what is possible in source code and on the processor. This is mostly the case for operations that don't translate directly to machine code. For example a conversion of a floating point variable to an integer is usually performed by code that is inserted inline or a library function call is introduced by the compiler at that point.

!! Compiler generated code should be thoroughly verified because it could include dead code or never taken code paths.

As iSYSTEM tools gather coverage information directly from the hardware, the analysis of compiler generated code isn't different as analysis from code written by the user. If source code is not available, the information is given on the object and assembler levels only.

Please check below example. See how the assignment statement "c=d;" uses code that isn't obviously part of the source. What is even more interesting is that this code includes conditionals and even function calls. In this particular case you see that part of the code wasn't executed, including a function call, because this is data dependent. This could lead to catastrophic failure if the function (that wasn't called) is never tested and has bugs.

```

void Type_Simple()
{
    otm(fn_Type_Simple);
    char c=0;
    unsigned char uc=0;
    int i=0;
    unsigned ui=0;
    long l=0;
    unsigned long ul=0;
    int iCount=0;
    long long ll=0x12345678;
    float f=(float)0.0;
    double d=1.583;

    for (iCount=-2;iCount<=2;++iCount)
    {
        c=iCount+3;
        d=c;
        c=d;
        uc=iCount;
        i=iCount;
        ui=iCount;
        l=iCount;
        ul=iCount;
        f=(float)iCount;
    }
}

```

		d=c;	
➔	20000444	881F0035	lbz r0,35(r31)
	20000448	7C000774	extsb r0,r0
	2000044C	7C030378	mr r3,r0
	20000450	480019F9	bl __floatsidf (20001E48)
	20000454	907F0038	stw r3,38(r31)
	20000458	909F003C	stw r4,3C(r31)
	2000045C	881F0035	lbz r0,35(r31)
	20000460	7C000774	extsb r0,r0
	20000464	2F800000	cmpi 7,0,r0,00
➔	20000468	409C0028	bc 04,1C,20000490
	2000046C	807F0038	lwz r3,38(r31)
	20000470	809F003C	lwz r4,3C(r31)
	20000474	3CA04070	lis r5,40700000
	20000478	38C00000	li r6,00
	2000047C	480018BD	bl __adddf3 (20001D38)
	20000480	7C691B78	mr r9,r3
	20000484	7C8A2378	mr r10,r4
	20000488	913F0038	stw r9,38(r31)
	2000048C	915F003C	stw r10,3C(r31)
	20000490	813F0038	lwz r9,38(r31)
	20000494	815F003C	lwz r10,3C(r31)
	20000498	913F0008	stw r9,08(r31)
	2000049C	915F000C	stw r10,0C(r31)



5. Compiler Optimizations

Another aspect is compiler optimizations. Compilers may

- merge several source code lines (e.g. merge string constant expressions, expression simplification,...),
- deem unnecessary code lines (e.g. dead code elimination, removal of unused symbols, common subexpression elimination,...) and/or
- reorder the execution flow (e.g. loop inversion, loop fusion,...).

!! When measuring code coverage, verify what source code lines generate object code and if not, check why. Any code change and recompilation may change the compiler optimization and could lead to untested code paths.

iSYSTEM tools take great care to show accurately what is actually going on. In iSYSTEM's winIDEA source code view every line that generates object code has a code coverage legend aligned to the left. Note that some “{“ and “}” are not creating object code and therefore have no code coverage legend aligned.

```
void Address_TestScopes()
{
  otm(fn_Address_TestScopes);
  int i=0,j=0;
  union uniA X;

  X.m_1=0x77;

  for (i=0;i<2;++i)
  {
    char c=4;
    X.m_1++;
    for (j=0;j<2;++j)
    {
      int X;
      ++c;
      X=c;
      if (c==1)
        ++X;
      else
        X+=2;
    }
  }

  ++X.m_1;
  ++iCounter;
  otm(fn_Address_TestScopesExit);
}
```

6. **Compiler Bugs**

Seldom, but still, compiler can have bugs. These bugs are rare on the code that is generated, as compilers are usually tested in this respect. Most of the time bugs are in the debug information generated by the compiler and used by other tools. It is important to remark that these happen seldom, but shouldn't be ruled out.

Conclusion

The requirements for code coverage analysis are constantly increasing in almost every market. Code coverage based on object code is a convenient method to test as close as possible to the final product.

Evaluating code coverage results keep in mind what we have discussed in this article. To achieve a high level of code coverage (some claim 100% is the right figure), review the object code and find test cases to cover the (not yet) tested code.

iSYSTEM tools provide code coverage analysis within the same tool that is used for code development. This way development costs and time are optimally used to detect and fix software bugs as soon as possible within the development cycle.

Consequently, choosing a processor or controller for your product should also include the requirement of providing on-chip trace capabilities.

Links

[Difference between Source and Object code Coverage](#)

[Offline versus Realtime Execution Coverage](#)

[Code Coverage: Relation between Source and Object Code](#)

References

[1] Beizer, Boris: *Software Testing Techniques*, Second edition, Von Nostrand Reinhold Company, Inc., 1990

[2] Hayhurst, K. J., Veerhusen D. S., Chilenski J. J., Rierson, L. K.: *NASA / TM-2001-210876 A Practical Tutorial on Modified Condition/Decision Coverage*, NASA Center for Aerospace Information (CASI)

[3] Certification Authorities Software Team (CAST): *Guidelines for Approving Source Code to Object Code Traceability*, Completed December 2002